

# Complexity of Target for pushdown RBN

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## Context

Reconfigurable Broadcast Networks (RBN) are a model for large groups of identical agents communicating via unreliable broadcast. A pushdown RBN (PRBN) is simply one where each agent is modelled by a pushdown transition system.

A PRBN is described by a tuple  $(Q, M, \Gamma, \Delta, q_0)$  where  $Q$  is a finite set of states,  $M$  a finite alphabet of messages,  $\Gamma$  a finite stack alphabet with a special letter  $\perp$ ,  $q_0$  an initial state, and

$$\Delta \subseteq Q \times \{\mathbf{br}(m), \mathbf{rec}(m) \mid m \in M\} \times \{\mathit{push}(\gamma), \mathit{pop}(\gamma), \mathit{nop} \mid \gamma \in \Gamma\} \times Q$$

a transition function.

A run starts with an arbitrarily large set of agents, all in state  $q_0$ . A step consists of one agent executing a transition with a broadcast  $\mathbf{br}(m)$ , and an arbitrary subset of other agents executing a transition with a matching reception  $\mathbf{rec}(m)$ .

Formally, a local configuration is a pair  $(q, \sigma) \in Q \times \Gamma^*$ . A local step  $(q, \sigma) \xrightarrow{\mathit{op}_1} (q', \sigma')$  is defined when there exists a transition  $(q, \mathit{op}_1, \mathit{op}_2, q')$  such that either:

- $\mathit{op}_2 = \mathit{push}(\gamma)$  and  $\sigma' = \gamma\sigma$
- $\mathit{op}_2 = \mathit{pop}(\gamma)$  and  $\sigma = \gamma\sigma'$
- $\mathit{op}_2 = \mathit{nop}$  and  $\sigma' = \sigma$

A configuration is a function from a finite non-empty set of agents  $\mathbb{A}$  to the set of local configurations  $C : \mathbb{A} \rightarrow Q \times \Gamma^*$ . An initial configuration is such that  $C(a) = (q_0, \perp)$  for all  $a \in \mathbb{A}$ . A step  $C \rightarrow C'$  is defined when there exists an agent  $a_{\mathbf{br}}$  such that there is a local step  $C(a_{\mathbf{br}}) \xrightarrow{\mathbf{br}(m)} C'(a_{\mathbf{br}})$  and for all  $a \neq a_{\mathbf{br}}$ , either

- there is a local step  $C(a) \xrightarrow{\mathbf{rec}(m)} C'(a)$ , or
- $C(a) = C'(a)$

A run is a sequence of consecutive steps, it is initial if it starts in an initial configuration.

Classical problems on such models ask whether a given set of configurations is reachable. A particular case of interest is the reachability of a configuration where all agents are in a given subset of states  $S \subseteq Q$ . This problem is called Target.

The open problem is to find tight complexity bounds on the Target problem for PRBN. The problem on finite-state RBN is known to be solvable in polynomial time. An easy reduction from Horn satisfiability shows PTIME-completeness.

On the other hand, one can show that the problem for PRBN is in NP: to witness Target, one can guess a set of messages  $M' \subseteq M$  of messages used, and two orders on this set, an order of appearance  $\leq_a$  and one of disappearance  $\leq_d$ .

It suffices to check that for all  $m \in M'$  there exists a run  $(q_0, \perp) \xrightarrow{*} (q, \sigma) \xrightarrow{*} (q_f, \sigma_f)$  only receiving messages smaller than  $m$  for  $\leq_a$  on the first part and messages from  $M'$  on the second part, as well as a run  $(q_0, \perp) \xrightarrow{*} (q', \sigma') \xrightarrow{*} (q_f, \sigma_f)$  only receiving messages from  $M'$  on the first part and from  $\{m' \in M' \mid m' <_d m\}$  on the second one.

This can be checked in polynomial time.

## Credits and related work

RBN were introduced (with finite-state systems) in [2], and the precise complexity of the main problems was analysed more finely in [3]. Pushdown RBN were introduced in [1], with in particular a proof that the coverability problem (whether there is a reachable configuration with at least one agent in a given state  $q_f$ ) is decidable in PTIME for this model.

## References

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- [3] Giorgio Delzanno et al. “On the Complexity of Parameterized Reachability in Reconfigurable Broadcast Networks”. In: *IARCS Annual Conference on Foundations of Software Technology and Theoretical Computer Science, FSTTCS 2012, December 15-17, 2012, Hyderabad, India*. Ed. by Deepak D’Souza, Telikepalli Kavitha, and Jaikumar Radhakrishnan. Vol. 18. LIPIcs. Schloss Dagstuhl - Leibniz-Zentrum für Informatik, 2012, pp. 289–300. DOI: 10.4230/LIPIcs.FSTTCS.2012.289. URL: <https://doi.org/10.4230/LIPIcs.FSTTCS.2012.289>.